



A MODIFIED CG-DESCENT AND DPR ALGORITHM FOR UNCONSTRAINED OPTIMIZATION

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Abstract: In this paper, a modified CG-DESCENT and DPR conjugate gradient algorithm is presented, which produces sufficient descent search direction at every iteration. This property depends neither on the line search used, nor on the convexity of the objective function. Moreover, we establish the global convergence for general nonlinear functions under suitable conditions. Numerical results show that the proposed method is more efficient than the previous ones.

Key words: Unconstrained optimization, conjugate gradient method, sufficient descent property, global convergence

Mathematics Subject Classification: 90C30, 65K05, 49M37

1 Introduction

Conjugate gradient methods comprise a class of unconstrained optimization algorithms which are characterized by low memory requirements and strong local and global convergence properties. For a general unconstrained optimization problem

$$\min f(x), \quad x \in \mathbb{R}^n, \tag{1.1}$$

where $f: \mathbb{R}^n \to \mathbb{R}$ is a continuously differentiable function. The iterative formula of the conjugate gradient method is given by

$$x_{k+1} = x_k + \alpha_k d_k, \quad k = 0, 1, \cdots,$$
(1.2)

where the step-length α_k is obtained by carrying out some line search, and the direction d_k is defined by

$$d_k = \begin{cases} -g_0, & \text{if } k = 0, \\ -g_k + \beta_k d_{k-1}, & \text{if } k > 0, \end{cases}$$
(1.3)

where β_k is a parameter. Some well-known formulas for β_k are the Fletcher-Reeves(FR) [2], Polak-Ribiere(PRP) [14], Liu-Storey(LS) [13], Dai-Yuan(DY) [4], the Conjugate Descent(CD) [8] and Hestenes-Stiefel(HS) [11] formulas, which are respectively given by

$$\beta_k^{DY} = \frac{\|g_k\|^2}{d_{k-1}^T y_{k-1}}, \quad \beta_k^{FR} = \frac{\|g_k\|^2}{\|g_{k-1}\|^2}, \quad \beta_k^{CD} = \frac{\|g_k\|^2}{-d_{k-1}^T g_{k-1}},$$

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$$\beta_k^{HS} = \frac{g_k^T y_{k-1}}{d_{k-1}^T y_{k-1}}, \quad \beta_k^{PR} = \frac{g_k^T y_{k-1}}{\|g_{k-1}\|^2}, \quad \beta_k^{LS} = \frac{g_k^T y_{k-1}}{-d_{k-1}^T g_{k-1}}$$

where $y_{k-1} = g_k - g_{k-1}$ and $\|\cdot\|$ stands for the Euclidean norm of vectors.

An important class of conjugate gradient algorithms is the hybrid conjugate gradient methods. Hu & Storey [12] and Dai & Yuan [5] proposed some hybrid methods which we call the H1 method and the H2 method, respectively, that is,

$$\beta_k^{H1} = \max\{0, \min\{\beta_k^{FR}, \beta_k^{PRP}\}\},$$
(1.4)

$$\beta_k^{H2} = \max\{0, \min\{\beta_k^{DY}, \beta_k^{HS}\}\}.$$
(1.5)

Gilbert and Nocedal [9] extended H1 to the case that

$$\beta_k^{H1} = \max\{-\beta_k^{FR}, \min\{\beta_k^{FR}, \beta_k^{PRP}\}\}.$$
(1.6)

Numerical performances show that the H1 and the H2 methods are better than the PRP method [5, 12].

Recently, there has been growing interest in the descent conjugate gradient methods. Hager and Zhang [10] proposed a new conjugate gradient method which was obtained by modifying the HS method and called CG-DESCENT method. The parameter β_k in the CG-DESCENT method is given by

$$\beta_k^N = \frac{g_k^T y_{k-1}}{d_{k-1}^T y_{k-1}} - 2 \frac{\|y_{k-1}\|^2 g_k^T d_{k-1}}{(d_{k-1}^T y_{k-1})^2},$$
(1.7)

$$\beta_k^{N+} = \max\{\beta_k^N, \eta_k\}, \eta_k = \frac{-1}{\|d_k\| \min\{\|g_k\|, \eta\}},$$
(1.8)

where $\eta > 0$ is a constant. Later, Yu and Guan [15] motivated by their work proposed a PRP conjugate gradient method as following

$$\beta_k^{DPR} = \frac{g_k^T y_{k-1}}{\|g_{k-1}\|^2} - C \frac{\|y_{k-1}\|^2 g_k^T d_{k-1}}{\|g_{k-1}\|^4}, \tag{1.9}$$

where parameter C essentially controls the relative weight between conjugant and descent. Zhang and Zhou [16] proposed two hybrid methods called NH1 and NH2 method as follows

$$NH1: d_k = -(1 + \beta_k^{\text{H1}} \frac{d_{k-1}^T g_k}{\|g_k\|^2})g_k + \beta_k^{\text{H1}} d_{k-1}, \qquad (1.10)$$

$$NH2: d_k = -(1 + \beta_k^{\text{H2}} \frac{d_{k-1}^T g_k}{\|g_k\|^2})g_k + \beta_k^{\text{H2}} d_{k-1}.$$
(1.11)

Obviously, these two new hybrid methods satisfy

$$g_k^T d_k = -\|g_k\|^2, (1.12)$$

which shows that they are descent and independent of any line search used. The global convergence of these two methods [16] are presented and numerical results also showed their efficiency in real computations.

In this paper, in order to obtain an efficiency method in real computations, based on the idea of the methods all above, we proposed a new method, which is a projection of the

CG-DESCENT and DPR conjugate algorithms. The search direction d_k has the following form

$$d_{k} = \begin{cases} -g_{0}, & \text{if } k = 0, \\ -(1 + \beta_{k}^{HZPR} \frac{d_{k-1}^{T}g_{k}}{\|g_{k}\|^{2}})g_{k} + \beta_{k}^{HZPR} d_{k-1}, & \text{if } k \ge 1, \end{cases}$$
(1.13)

where

$$\beta_k^{HZPR} = \max\{0, \min\{\beta_k^N, \beta_k^{DPR}\}\}.$$
(1.14)

From the method above, we can easily obtain (1.12). For convenience, we call the method above as the HZPR method.

The rest of this paper are organized as follows. In the next section, we prove the global convergence of the method (1.13) for general nonlinear functions with strong Wolfe line search. In section 3, we report some numerical results to test the proposed method.

2 Algorithm and Convergence Analysis

First, we make the following standard assumptions for the objective function, which have been used often in the literature to analyze the global convergence of conjugate methods with inexact line search.

Assumption (A)

- (H1) The level set $\Omega = \{x \in \mathbb{R}^n | f(x) \le f(x_0)\}$ is bounded.
- (H2) In some neighborhood N of Ω , f is continuously differentiable and its gradient is Lipschitz continuous, namely, there exists a constant L > 0 such that

$$\|g(x) - g(y)\| \le L \|x - y\|, \quad \forall x, y \in N.$$

Assumption (A) implies that there exists a positive constant $\hat{\gamma}$ and B such that

$$\|g(x)\| \le \widehat{\gamma}, \quad \forall x \in \Omega, \tag{2.1}$$

and

$$\|x - y\| \le B, \quad \forall x, y \in \Omega.$$

$$(2.2)$$

Now we introduce the steps of the HZPR algorithm as following.

Algorithm 2.1 (HZPR Method).

Step 0: Choose an initial point $x_0 \in \mathbb{R}^n$. Let k = 0.

Step 1: Compute d_k by Eq. (1.13), where β_k^{HZPR} is computed by (1.14).

Step 2: Determine α_k by the strong Wolfe line search

$$f(x_k + \alpha_k d_k) - f(x_k) \le \delta \alpha_k g_k^T d_k, \tag{2.3}$$

$$|g(x_k + \alpha_k d_k)^T d_k| \le \sigma |g_k^T d_k|, \qquad (2.4)$$

where $0 < \delta < \sigma < 1$.

Step 3: Let $x_{k+1} = x_k + \alpha_k d_k$.

Step 4: Let k := k + 1 and go to step 1.

The following well-known lemma, called the zoutendijk condition, which was originally given in [18].

Lemma 2.2. Suppose Assumption (A) holds. Consider any method in the form (1.2), where d_k is a descent direction and α_k satisfies the strong Wolfe condition (2.3) and (2.4). Then we have

$$\sum_{k=0}^{\infty} \frac{(d_k^T g_k)^2}{\|d_k\|^2} < \infty.$$
(2.5)

We now establish the global convergence theorem for the HZPR algorithm in a similar way to Theorem 3.1 in [6]. First, we give a useful lemma about β_k^{HZPR} in (1.14) which plays an important role in the global convergence analysis.

Lemma 2.3. Suppose Assumption (A) holds. $\{x_k\}$ is generated by algorithm 2.1. If there exist a constant $\epsilon > 0$ such that $||g_k|| > \epsilon$ for all k > 0, then there exist a positive constant D such that

$$|\beta_k^{HZPR}| \le D \|s_{k-1}\|, \tag{2.6}$$

where $s_{k-1} = x_k - x_{k-1}$.

Proof. From (2.4) and (1.12), we have

$$\begin{aligned} |d_{k-1}^T y_{k-1}| &= |d_{k-1}^T g_k - d_{k-1}^T g_{k-1}| \ge |d_{k-1}^T g_{k-1}| - |d_{k-1}^T g_k| \\ &\ge |d_{k-1}^T g_{k-1}| - \sigma |g_{k-1}^T d_{k-1}| = (1-\sigma) ||g_{k-1}||^2. \end{aligned}$$

The above inequality together with (1.12) and (2.4), we have

$$\begin{split} |\beta_{k}^{HZPR}| &= |\max\{0, \min\{\beta_{k}^{DPR}, \beta_{k}^{N}\}| \\ &\leq \max\{\frac{\|g_{k}\|\|y_{k-1}\|}{\|g_{k-1}\|^{2}} + \frac{C\|y_{k-1}\|^{2}\sigma|g_{k-1}^{T}d_{k-1}|}{\|g_{k-1}\|^{4}}, \\ & \frac{\|g_{k}\|\|y_{k-1}\|}{|d_{k-1}^{T}y_{k-1}|} + \frac{2\|y_{k-1}\|^{2}\sigma|g_{k-1}^{T}d_{k-1}|}{(d_{k-1}^{T}y_{k-1})^{2}}\} \\ &\leq \max\{\frac{\|g_{k}\|\|y_{k-1}\|}{\|g_{k-1}\|^{2}} + \frac{C\|y_{k-1}\|^{2}\sigma\|g_{k-1}\|^{2}}{\|g_{k-1}\|^{4}}, \\ & \frac{\|g_{k}\|\|y_{k-1}\|}{(1-\sigma)\|g_{k-1}\|^{2}} + \frac{2\|y_{k-1}\|^{2}\sigma\|g_{k-1}\|^{2}}{(1-\sigma)\|g_{k-1}\|^{4}}\} \\ &\leq \max\{\frac{L\widehat{\gamma} + C\sigma L^{2}B}{\epsilon^{2}}, \frac{L\widehat{\gamma} + 2\sigma L^{2}B}{(1-\sigma)\epsilon^{2}}\}\|s_{k-1}\|. \end{split}$$

Defining

$$D = \max\{\frac{L\widehat{\gamma} + C\sigma L^2 B}{\epsilon^2}, \frac{L\widehat{\gamma} + 2\sigma L^2 B}{(1-\sigma)\epsilon^2}\}$$

then we have the result (2.6). The proof is complete.

The next lemma corresponds to Lemma 3.4 in [3] and Lemma 2.3 in [6].

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Lemma 2.4. Suppose Assumption (A) holds. $\{x_k\}$ is generated by algorithm 2.1. If there exist a constant $\epsilon > 0$ such that $||g_k|| > \epsilon$ for all k > 0, then we have

$$\sum_{k=0}^{\infty} \|u_k - u_{k-1}\|^2 < \infty,$$
(2.7)

where $u_k = \frac{d_k}{\|d_k\|}$.

Proof. Since $d_k \neq 0$ follows from (1.12) and $||g_k|| > \epsilon$, so u_k is well-defined. We rewrite (1.13) as following

$$d_{k} = -(1 + \beta_{k}^{HZPR} \frac{d_{k-1}^{T}g_{k}}{\|g_{k}\|^{2}})g_{k} + \beta_{k}^{HZPR} d_{k-1} = v_{k} + \beta_{k}^{HZPR} d_{k-1}.$$
 (2.8)

By defining

So, we have

$$r_{k} = \frac{v_{k}}{\|d_{k}\|}, \quad \delta_{k} = \frac{\beta_{k}^{HZPR} \|d_{k-1}\|}{\|d_{k}\|}.$$
$$u_{k} = r_{k} + \delta_{k} u_{k-1}.$$
(2.9)

Then we have from the fact that $||u_k|| = ||u_{k-1}|| = 1$

$$||r_k|| = ||u_k - \delta_k u_{k-1}|| = ||u_{k-1} - \delta_k u_k||.$$
(2.10)

Using the condition $\delta_k \ge 0$, the triangle inequality and (2.10), we have

$$||u_k - u_{k-1}|| \le ||(1 + \delta_k)(u_k - u_{k-1})|| \le ||u_k - \delta_k u_{k-1}|| + ||u_{k-1} - \delta_k u_k|| \le 2||r_k||.$$
(2.11)

From (2.1), (2.4) and (2.6), we have

$$|\beta_k^{HZPR}| \frac{|g_k^T d_{k-1}|}{\|g_k\|^2} \le D \|s_{k-1}\| \frac{\sigma |g_{k-1}^T d_{k-1}|}{\|g_k\|^2} \le D \frac{\sigma \widehat{\gamma}^2}{\epsilon^2} \|s_{k-1}\| \doteq M \|s_{k-1}\|.$$
(2.12)

From (2.2), (2.12) and (2.8), there exist a constant $M_1 \ge 0$ such that

$$||v_k|| \le ||g_k|| + M||s_{k-1}|| ||g_k|| \le M_1.$$
(2.13)

From the definition of r_k , (2.6) and (2.13), we have

$$\sum_{k=0}^{\infty} \|r_k\|^2 = \sum_{k=0}^{\infty} \frac{\|v_k\|^2}{\|d_k\|^2} \le \sum_{k=0}^{\infty} \frac{M_1^2}{\|d_k\|^2} = \sum_{k=0}^{\infty} \frac{M_1^2}{\|g_k\|^4} \frac{\|g_k\|^4}{\|d_k\|^2} \le \frac{M_1^2}{\epsilon^2} \sum_{k=0}^{\infty} \frac{\|g_k\|^4}{\|d_k\|^2} < \infty.$$
(2.14)

Together with (2.11), we have (2.7).

The next theorem establishes the global convergence of the HZPR method. The proof of it is similar to Theorem 4.3 in [9] and Theorem 3.1 in [6].

Theorem 2.5. Suppose Assumption (A) holds. $\{x_k\}$ is generated by algorithm 2.1. Then we have

$$\liminf_{k \to \infty} \|g_k\| = 0. \tag{2.15}$$

Proof. Assume that the conclusion (2.15) is not true, then there exist a constant $\varepsilon > 0$ such that for all k, $||g_k|| > \varepsilon$. From the definition of u_k , we observe that for any $l \ge k$,

$$x_{l} - x_{k} = \sum_{j=k}^{l-1} (x_{j+1} - x_{j}) = \sum_{j=k}^{l-1} \|s_{j}\| u_{k} + \sum_{j=k}^{l-1} \|s_{j}\| (u_{j} - u_{k}).$$
(2.16)

By the triangle inequality, from the fact $||u_k|| = 1$, so we have

$$\sum_{j=k}^{l-1} \|s_j\| \le \|x_l - x_k\| + \sum_{j=k}^{l-1} \|s_j\| \|u_j - u_k\| \le B + \sum_{j=k}^{l-1} \|s_j\| \|u_j - u_k\|.$$
(2.17)

Let Δ be a positive integer, chosen large enough that

$$\Delta \ge 4BD,\tag{2.18}$$

where B and D appear in (2.2) and (2.6). By lemma 2.4, we can find a large enough k_0 that

$$\sum_{k \ge k_0} \|u_k - u_{k-1}\|^2 \le \frac{1}{4\Delta}.$$
(2.19)

If $j > k > k_0$ and $j - k \le \Delta$, then by (2.19) and Cauchy-Schwarz inequality, we have

$$\begin{aligned} \|u_{j} - u_{k}\| &\leq \sum_{i=k}^{j-1} \|u_{i+1} - u_{i}\|^{2} \\ &\leq \sqrt{j-k} (\sum_{i=k}^{j-1} \|u_{i+1} - u_{i}\|^{2})^{1/2} \\ &\leq \sqrt{\Delta} (\frac{1}{4\Delta})^{1/2} = \frac{1}{2}. \end{aligned}$$

Combining this with (2.2) and (2.17) yields

$$\sum_{j=k}^{l-1} \|s_j\| \le 2B. \tag{2.20}$$

From (2.6), (2.8) and (2.13) we have

$$||d_l||^2 \le (||v_k|| + |\beta_l^{HZPR}|||d_{l-1}||)^2 \le 2M_1^2 + 2D^2 ||s_{l-1}||^2 ||d_{l-1}||^2.$$
(2.21)

Defining $S_i = 2D^2 ||s_i||^2$, by induction, we obtain

$$\begin{aligned} \|d_{l}\|^{2} &\leq 2M_{1}^{2} + S_{l-1} \|d_{l-1}\|^{2} \\ &\leq 2M_{1}^{2} (1 + S_{l-1} + S_{l-1}S_{l-2} + \dots + S_{l-1}S_{l-2} \dots S_{k_{0}+1}) \\ &+ \|d_{k_{0}}\|^{2} S_{l-1}S_{l-2} \dots S_{k_{0}}. \end{aligned}$$

Then we have

$$\|d_l\|^2 \le \begin{cases} 2M_1^2 + S_{k_0} \|d_{k_0}\|^2, & \text{if } l = k_0 + 1, \\ 2M_1^2 (1 + \sum_{i=k_0+1}^{l-1} \prod_{j=i}^{l-1} S_j) + \|d_{k_0}\|^2 \prod_{j=k_0}^{l-1} S_j, & \text{if } l > k_0 + 1. \end{cases}$$
(2.22)

Let us consider as follows a product of Δ consecutive S_i , where $k \geq k_0$,

$$\prod_{j=k}^{k+\Delta-1} S_j = \prod_{j=k}^{k+\Delta-1} 2D^2 \|s_j\|^2 = (\prod_{j=k}^{k+\Delta-1} \sqrt{2}D \|s_j\|)^2$$
$$\leq (\frac{\sum_{j=k}^{k+\Delta-1} \sqrt{2}D \|s_j\|}{\Delta})^{2\Delta} \leq (\frac{2\sqrt{2}BD}{\Delta})^{2\Delta} \leq \frac{1}{2^{\Delta}}$$

The product of Δ consecutive S_j is bounded by $\frac{1}{2^{\Delta}}$, it follows that the sum in (2.22) is bounded, and the bound is independent of l. This bound for $||d_l||$, independent of $l > k_0$, contradicts (2.5), hence we have (2.15). The proof is complete.

3 Numerical Experiments

In this section, we do some numerical experiments to test the performance of the HZPR method and compare it with some existing methods for solving large scale unconstrained optimization problems. All codes are written in Fortran and ran on IBM T60 PC with two 1.83 GHz CPU and 2.5GB RAM.

The test problems are the unconstrained problems from Neculai Andrei [1]. For each problem, the dimension n is set to 1 000 and 10 000. The parameters in the strong wolfe conditions are as follows: $\sigma = 0.9$ and $\delta = 0.1$, and C = 1 in (1.14). We stop the iteration if the inequality $||g_k|| \leq 10^{-6}$ is satisfied.

We compare the performances of the HZPR method with that of the CG-DESCENT method [10] and the MPRP method [17]. The CG-DESCENT codes can be obtained from Hager's page at http://www.math.ufl.edu/hager/papers/CG.

Table 1 lists the results of the HZPR method, the CG-DESCENT method and the MPRP method which gives the total number of iterations(iter), the total number of function evaluations(fn), the total number of gradient evaluations(gn) and the cpu time(time) in seconds.

We adopt the performance profiles by Dolan and More [7] to compare the performance among the tested methods. That is, for each method, we plot the fraction P of problems for which the method is within a factor τ of the best time. The left side of the figure gives the percentage of the test problems for which a method is the fastest; the right side gives the percentage of the test problems that are successfully solved by each of the methods. The top curve is the method that solved the most problems in a time that are within a factor τ of the best time. Figure 1-4 are the performance profile measured by CPU time, the number of iterations, the number of function evaluations and the number of gradient evaluations, respectively.

Table 1: The result of HZPR, MPRP and CG-DESCENT

Problem	N	CG-DESCENT	MPRP	HZPR
		iter/fn/gn/time	iter/fn/gn/time	iter/fn/gn/time
ROTH	1000	14/31/22/0.00E+00	12/30/22/1.56E-02	12/29/23/1.56E-02
ROTH	10000	18/39/27/9.38E-02	15/36/27/9.38E-02	11/25/19/4.69E-02
TRIGMETRIC	1000	100/211/117/2.34E-01	73/151/81/1.56E-01	70/146/78/1.56E-01
TRIGMETRIC	10000	94/196/109/2.17E+00	79/164/89/1.81E+00	86/183/104/2.06E+00
ROSENBROCK	1000	36/120/94/1.56E-02	40/104/76/1.56E-02	40/116/90/3.12E-02
ROSENBROCK	10000	34/111/87/1.56E-01	38/98/70/1.56E-01	46/132/102/2.03E-01
WHITEHOLST	1000	39/117/87/1.56E-02	40/109/77/1.56E-02	43/132/101/1.56E-02
WHITEHOLST	10000	39/123/94/1.88E-01	41/110/76/1.88E-01	42/134/105/2.19E-01
BEALEU63	1000	17/36/21/1.56E-02	12/25/14/1.56E-02	17/35/21/1.56E-02
BEALEU63	10000	17/36/21/7.81E-02	12/25/14/4.69E-02	17/35/21/7.81E-02
PENALTY	1000	36/70/42/1.56E-02	36/69/41/1.56E-02	33/66/39/1.56E-02
PENALTY	10000	42/80/48/1.56E-01	38/76/43/1.41E-01	38/76/42/1.41E-01
PQUADRATIC	1000	188/377/189/7.81E-02	188/377/189/6.25E-02	188/377/189/6.25E-02
PQUADRATIC	10000	598/1197/599/2.20E+00	598/1197/599/ 2.28E+00	598/1197/599/2.31E+00
RAYDAN1	1000	252/381/377/2.66E-01	242/365/363/2.34E-01	232/352/346/2.34E-01

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		Table 1 – continued from previous page						
Problem	N	CG-DESCENT iter/fn/gn/time	MPRP iter/fn/gn/time	HZPR iter/fn/gn/time				
RAYDAN1	10000	829/1101/1388/7.98E+00	839/1101/1418/8.17E+00	808/1061/1365/7.89E+00				
RAYDAN2	1000	5/11/7/0.00E+00	6/13/9/0.00E+00	6/12/8/1.56E-02				
RAYDAN2	10000	6/13/9/6.25E-02	6/13/9/6.25E-02	6/13/9/7.81E-02				
DIAGONAL1	1000	307/430/493/3.59E-01	330/463/529/3.91E-01	296/416/474/3.59E-01				
DIAGONALI DIAGONAL2	10000	989/1218/1753/1.15E+01 201/382/247/2 50E=01	998/1235/1761/1.18E+01 233/436/307/2 97E-01	991/1216/1761/1.17E+01 207/386/273/2 50E=01				
DIAGONAL2	10000	615/1175/809/7.88E+00	641/1161/873/8.17E+00	616/1174/841/8.03E+00				
DIAGONAL3	1000	262/374/414/4.22E-01	270/385/427/4.38E-01	261/373/412/4.22E-01				
DIAGONAL3	10000	889/1095/1574/1.42E+01	881/1095/1550/1.42E+01	885/1095/1562/1.42E+01				
HAGER	1000	50/85/77/7.81E-02 02/144/151/1.22E+00	50/85/77/7.81E-02 02/142/140/1.21E+00	50/84/78/7.81E-02 02/142/152/1 22E 00				
GTRIDIAG1	10000	26/45/37/0.00E+00	26/44/36/1.56E-02	26/44/36/0.00E+00				
GTRIDIAG1	10000	26/43/37/1.25E-01	26/42/38/1.41E-01	26/42/38/1.25E-01				
TRIDIAG1	1000	18/37/21/0.00E+00	21/43/29/1.56E-02	22/45/27/1.56E-02				
TRIDIAG1	10000	18/37/21/7.81E-02	21/43/29/9.38E-02	24/49/27/9.38E-02				
TETERMS	10000	9/19/13/1.88E-01	12/26/14/1.56E-02 12/26/18/2.66E-01	9/20/14/2.03E-02				
GTRIDIAG2	1000	50/92/60/1.56E-02	47/85/58/3.12E-02	46/84/56/1.56E-02				
GTRIDIAG2	10000	52/95/63/2.97E-01	54/99/65/3.12E-01	53/100/61/2.97E-01				
DIAGONAL4	1000	4/9/6/0.00E+00	4/9/6/0.00E+00	4/9/6/0.00E+00				
DIAGONAL5	10000	3/8/5/1 56E=02	4/9/0/1.50E-02 3/8/5/1.56E-02	$\frac{4}{9}/\frac{6}{1.30E} = 02$ $\frac{3}{8}/\frac{5}{0.00E} = 00$				
DIAGONAL5	10000	3/8/5/1.09E-01	3/8/5/1.09E-01	3/8/5/1.09E-01				
HIMMELB	1000	9/22/14/0.00E+00	9/22/14/0.00E+00	9/22/14/0.00E+00				
HIMMELB	10000	9/22/14/4.69E-02	9/22/14/4.69E-02	9/22/14/3.12E-02				
GPSC1 CPSC1	1000	726/1152/1471/1.06E+00 840/1224/1707/1.17E+01	488/845/854/ 7.34E-01	564/898/1124/8.28E-01				
PSC1	1000	15/29/18/3.12E-02	12/24/15/1.56E-02	13/26/17/3.12E-02				
PSC1	10000	12/26/15/1.88E-01	11/23/13/1.56E-01	12/24/15/1.56E-01				
POWELL	1000	116/236/135/4.69E-02	216/434/235/6.25E-02	75/153/93/3.12E-02				
POWELL BD1	10000	581/1186/659/2.00E+00 25/64/56/3 12E-02	178/358/197/6.41E-01 17/47/42/1 56E-02	328/660/393/1.19E+00 24/64/53/3 12E-02				
BD1 BD1	10000	25/64/56/2.97E-01	17/47/42/2.34E-01	24/64/53/2.97E-01				
MARATOS	1000	52/159/127/1.56E-02	70/230/186/3.12E-02	70/233/193/3.12E-02				
MARATOS	10000	51/168/134/2.34E-01	67/198/158/2.97E-01	74/235/193/3.44E-01				
QDP ODP	1000	135/271/162/4.69E-02 $442/885/544/1.91E\pm00$	137/275/165/6.25E-02 $430/861/523/1.89E\pm00$	136/273/163/6.25E-02 $430/861/525/1.91E\pm00$				
WOOD	10000	189/416/240/7.81E-02	269/594/339/1.09E-01	146/344/219/6.25E-02				
WOOD	10000	183/421/256/7.50E-01	248/528/308/1.00E+00	142/338/218/6.09E-01				
HIEBERT	1000	79/257/197/4.69E-02	76/248/198/3.12E-02	100/368/306/6.25E-02				
OF1	10000	1/236/190/3.28E-01 189/379/190/6.25E-02	189/379/190/7 81E-02	98/327/264/4.69E-01 189/379/190/6 25E-02				
QF1	10000	600/1201/601/1.95E+00	600/1201/601/2.03E+00	600/1201/601/2.03E+00				
QP1	1000	15/29/18/0.00E + 00	14/29/17/0.00E+00	14/30/17/0.00E + 00				
QP1 OP2	10000	17/35/21/7.81E-02 42/122/100/0.28E.02	16/35/20/7.81E-02 55/162/124/1.00E-01	17/36/21/7.81E-02				
QP2	10000	39/133/103/9.06E-01	50/152/116/1.05E+00	43/136/105/9.38E-01				
QF2	1000	393/687/501/1.25E-01	394/688/503/1.41E-01	390/683/496/1.41E-01				
QF2	10000	1253/2167/1601/4.39E+00	1276/2191/1646/4.62E+00	1251/2164/1598/4.55E+00				
EP1	10000	3/7/5/3.12E-02	3/7/5/3.12E-02	3/7/5/1.56E-02				
TRIDIAG2	1000	39/63/56/1.56E-02	37/61/52/1.56E-02	38/62/54/1.56E-02				
TRIDIAG2	10000	42/68/65/1.72E-01	39/61/62/1.72E-01	37/60/60/1.72E-01				
TRIDIA	1000	356/713/357/1 41E-01	358/717/359/1 41E-01	358/717/359/1 41E=01				
TRIDIA	10000	1175/2351/1176/4.45E+00	1176/2353/1177/4.59E+00	1177/2355/1178/4.61E+00				
ARWHEAD	1000	12/28/19/1.56E-02	14/37/27/0.00E+00	9/22/16/1.56E-02				
ARWHEAD	10000	10/22/15/4.69E-02 12/20/21/1 56E 02	8/19/13/4.69E-02 10/22/15/0.00E+00	9/24/18/4.69E-02				
NONDIA	10000	11/37/32/6.25E-02	9/29/22/4.69E-02	11/30/23/4.69E-02				
NONDQUAR	1000	$16252/32520/17640/7.30E \pm 00$	12041/24093/12529/5.47E+00	6501/13096/8345/3.17E+00				
DQDRTIC	1000	7/15/8/0.00E+00	7/15/8/0.00E+00	7/15/8/0.00E+00				
EG2	10000	10/21/11/4.09E-02 125/285/229/1.72E-01	124/257/238/1.56E-01	35/99/96/6.25E-02				
EG2	10000	1979/3584/5207/2.95E+01	567/1172/1672/9.42E+00	141/353/261/1.89E+00				
DIXMAANA	1000	9/19/10/1.56E-02	8/17/9/1.56E-02	9/19/10/0.00E+00				
DIXMAANA	10000	9/19/10/9.38E-02 21/55/35/3 12E 02	7/15/8/7.81E-02 20/54/34/3 12E 02	8/17/9/9.38E-02 22/58/38/1 56E 02				
DIXMAANB	10000	22/57/38/2.97E-01	22/58/38/2.97E-01	22/59/39/2.97E-01				
DIXMAANC	1000	30/81/54/4.69E-02	26/70/45/3.12E-02	26/70/46/3.12E-02				
DIXMAANC	10000	31/83/55/4.06E-01	27/72/48/3.59E-01	24/66/43/3.28E-01				
DIXMAANE	1000	176/339/191/1.88E-01 $465/917/486/4.69E\pm00$	162/319/169/1.72E-01 $451/902/458/4.56E\pm00$	171/328/187/1.72E-01 $468/925/494/4.80E\pm00$				
PPQ	10000	159/319/160/2.58E+00	160/321/161/2.61E+00	159/319/160/2.59E+00				
PPQ	10000	28/57/33/5.21E+01	29/59/34/5.37E+01	29/59/34/5.37E+01				
BT	1000	47/96/49/1.56E-02 27/75/28/1.72E.01	47/95/48/1.56E-02	40/81/41/1.56E-02 41/82/42/2.02E_01				
APQ	10000	189/379/190/6.25E-02	189/379/190/6.25E-02	189/379/190/6.25E-02				
APQ	10000	600/1201/601/1.95E+00	600/1201/601/2.05E+00	600/1201/601/2.05E+00				
TPQ	1000	177/355/178/7.81E-02	177/355/178/9.38E-02	177/355/178/7.81E-02				
EDENSCH	10000	29/49/42/1.56E-02	27/46/39/1.56E-02	27/46/38/1.56E-02				
EDENSCH	10000	27/47/42/1.41E-01	26/43/38/1.41E-01	26/44/40/1.41E-01				
VARDIM	1000	36/74/38/1.56E-02	36/73/38/1.56E-02	38/78/41/1.56E-02				
VARDIM S1	10000	40/93/47/1.72E-01 2000/4001/2002/7 19E-01	40/93/47/1.88E-01 2000/4001/2002/7 50E-01	40/93/47/1.88E-01 1999/3999/2001/7 50E-01				
LIARWHD	1000	20/42/27/0.00E+00	24/52/37/1.56E-02	24/53/36/0.00E+00				
LIARWHD	10000	26/62/43/1.25E-01	25/56/36/1.09E-01	25/55/36/1.25E-01				
DIAGONAL6	1000	5/11/6/4.69E-02	5/11/6/4.69E-02	5/11/6/6.25E-02				
DIXON3DQ	1000	1989/3979/1991/7.03E-01	1993/3987/1995/7.19E-01	1993/3987/1995/7.03E-01				
DIXMAANF	1000	227/484/295/2.50E-01	338/680/401/3.59E-01	206/436/260/2.34E-01				
DIAMAANF	10000	097/1402/903/7.69E+00 234/478/270/2.66E-01	1364/2912/1014/1.51E+01 275/567/304/3 28E-01	227/464/255/2 50E-01				
		. ,,,		Continued on next page				

Table 1 – continued from previous page					
Problem	N CG-DESCENT MPRP		HZPR		
		iter/fn/gn/time	iter/fn/gn/time	iter/fn/gn/time	
DIXMAANG	10000	568/1129/631/6.17E+00	1186/2396/1247/1.28E+01	691/1371/772/7.61E+00	
DIXMAANH	1000	259/562/343/3.12E-01	35/145/124/7.81E-02	240/535/332/2.97E-01	
DIXMAANH	10000	61/206/184/1.22E+00	61/195/174/1.16E+00	758/1576/957/8.47E+00	
DIXMAANI	1000	157/313/160/1.56E-01	160/317/165/1.56E-01	157/314/159/1.72E-01	
DIXMAANI	10000	469/919/497/4.78E+00	444/891/448/4.53E+00	458/912/474/4.70E+00	
DIXMAANJ	1000	215/454/277/2.50E-01	259/522/307/2.81E-01	245/509/304/2.66E-01	
DIXMAANJ	10000	699/1468/946/7.97E+00	827/1699/934/8.81E+00	650/1354/842/7.34E+00	
DIXMAANK	1000	231/486/284/2.81E-01	212/456/264/2.66E-01	233/497/298/2.81E-01	
DIXMAANK	10000	653/1329/770/7.59E+00	1593/3180/1735/1.80E+01	612/1258/746/7.31E+00	
DIXMAANL	1000	5903/11878/7483/8.34E+00	14912/31009/16324/5.14E+01	5002/10585/5723/6.97E+00	
ENGVAL1	1000	29/52/39/1.56E-02	30/53/41/1.56E-02	27/48/36/1.56E-02	
ENGVAL1	10000	26/44/38/1.25E-01	28/50/40/1.41E-01	25/45/36/1.25E-01	
FLETCHCB	1000	2946/6017/3073/1.33E+00	$2926/6022/3100/1.36E\pm00$	$2930/6037/3108/1.38E\pm00$	
COSINE	1000	13/29/26/3 $12E=02$	11/26/23/1.56E=02	11/27/22/1 56E=02	
COSINE	10000	13/33/30/1 88E-01	12/28/27/1 72E-01	12/28/26/1.56E=01	
DENSCHNB	1000	$8/17/9/0.00E\pm00$	$6/13/7/0.00E\pm00$	8/17/9/1 56E-02	
DENSCHNB	10000	8/17/9/3 12E-02	6/13/7/1 56E-02	8/17/9/3 12E-02	
DENSCHNE	1000	30/71/56/3 12E-02	21/50/40/1 56E-02	22/53/43/1 56E=02	
DENSCHNE	10000	26/63/52/1 56E-01	21/50/40/1 25E-01	22/53/43/1 41E-01	
SINOUAD	10000	551/1223/761/7 97E-01	$1/83/3175/1880/2.05E\pm00$	290/772/570/5 47E-01	
SINOUAD	10000	2202 / 4762 / 2770 / 2 02E 01	2082/8502/4002/5 41E 01	2745/0250/6528/6 44E + 01	
BICCSP1	10000	500/1001/501/1 72E 01	500/1001/501/1 72E 01	500/1001/501/1 88E 01	
BIGGSB1	10000	5000/1001/5001/1.72E-01	5001/1001/501/1.72E-01	5000/1001/5001/1.88E-01	
PPO2	10000	0/1/1/1 56E 02	$0/1/1/0.00 \pm 0.00$	0/1/1/1 56E 02	
PPO2	10000	0/1/1/1.30E = 02 0/1/1/1.17E + 00	$0/1/1/0.00E \pm 00$	0/1/1/1.50E=02 0/1/1/1.16E + 00	
502	10000	52/107/54/1 56E 02	52/107/54/2 19E 02	52/107/54/1 56E 02	
5Q2	1000	177/255/178/5 78D 01	177/255/178/6 00E 01	177/255/178/5 04E 01	
GENDOGE	10000	1/1/303/1/8/3./8E-01	177/333/178/0.09E-01	177/355/178/5.94E-01	
GENROSE	1000	0457/15508/7001/2.80E+00	4238/8498/4290/1.84E+00	100 (000 (140 (4 COE 00	
NONDIA	1000	3336/7054/4709/1.42E+00	3041/6327/4693/1.36E+00	109/232/143/4.69E-02	
PENALTYI	1000	14/29/15/0.00E+00	14/29/15/0.00E+00	14/29/15/0.00E+00	
PENALTYI	10000	16/33/17/7.81E-02	16/33/17/6.25E-02	16/33/17/6.25E-02	
POWER	1000	10456/20913/10457/3.62E+00	10454/20909/10455/3.77E+00	10457/20915/10458/3.75E+00	
FREUOTH	1000	120/233/157/7.81E-02	69/130/105/4.69E-02	48/100/76/3.12E-02	
FREUOTH	10000	114/207/185/8.28E-01	52/104/89/3.91E-01	53/101/94/4.22E-01	
SROSENBR	1000	32/96/72/1.56E-02	36/88/62/1.56E-02	42/129/100/1.56E-02	
SROSENBR	10000	36/101/77/1.72E-01	38/92/65/1.72E-01	43/138/111/2.34E-01	
WOODS	1000	369/784/428/1.56E-01	332/711/394/1.41E-01	309/659/364/1.41E-01	
WOODS	10000	211/466/272/9.06E-01	263/564/311/1.14E+00	138/321/200/6.41E-01	
DQRTIC	1000	31/63/32/1.56E-02	31/63/32/1.56E-02	31/63/32/1.56E-02	
DQRTIC	10000	37/75/38/1.41E-01	36/73/37/1.41E-01	36/73/37/1.41E-01	
NONCVXU2	1000	2529/4386/3203/4.14E+00	2449/4237/3114/4.05E+00	2661/4408/3579/4.52E+00	
BROYDN7D	1000	1007/2304/1333/8.75E-01	985/2321/ 1366/8.91E-01	544/1290/768/5.00E-01	
BROWNAL	1000	6/19/19/0.00E+00	6/17/17/0.00E+00	6/17/17/0.00E+00	
BROWNAL	10000	7/22/22/9.38E-02	7/19/19/9.38E-02	7/19/19/9.38E-02	
GENHUMPS	1000	778/1658/892/2.11E+00	654/1407/764/ 1.94E+00	655/1438/805/1.89E+00	
BDEXP	1000	14/29/15/4.69E-02	14/29/15/4.69E-02	14/29/15/4.69E-02	
BDEXP	10000	14/29/15/3.59E-01	14/29/15/3.91E-01	14/29/15/3.75E-01	



Figure 1: Performance profiles based on CPU time

From the results of our numerical experiments, we can see that the HZPR method performs better than the CG-DESCENT method and the MPRP method, which implies



Figure 2: Performance profiles based on iterations



Figure 3: Performance profiles for the number of function evaluations



Figure 4: Performance profiles for the number of gradient evaluations

that the HZPR method is efficient in real computation.

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