



## AN ELEMENTARY PROOF OF CONVERGENCE FOR THE FORWARD-BACKWARD SPLITTING ALGORITHM

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**ABSTRACT.** The forward-backward splitting (FBS) algorithm is a quite general iterative method that includes, as particular cases, the projected gradient descent algorithm for constrained minimization, the CQ algorithm for the split feasibility problem, the projected Landweber algorithm for constrained least squares, and the simultaneous orthogonal projection algorithm for the convex feasibility problem. The FBS algorithm involves iterating with respect to an averaged operator that is the product of two firmly non-expansive operators, one of which is Moreau's proximity operator. The usual proof of convergence employs the Krasnosel'skii-Mann Theorem. This proof depends, therefore, on knowing that if the gradient of a convex differentiable function is non-expansive, then it is firmly non-expansive, and that the composition of averaged operators is again averaged. Neither of these results is trivial to prove, especially the former. In this paper we give an elementary proof of convergence of the FBS algorithm that does not rely on these results.

### 1. BACKGROUND

The usual proof of convergence of the forward-backward splitting algorithm, as presented, for example, by Combettes and Wajs [11], relies on properties of averaged and firmly non-expansive operators, and on the Krasnosel'skii-Mann (KM) Theorem [16, 18]. In this section we elaborate on these points. Details can be found in [14, 1, 6] and in the references therein.

**1.1. Firmly Non-expansive and Averaged Operators.** In this subsection we review the basic theory of firmly non-expansive and averaged operators.

**Definition 1.1.** An operator  $T : \mathbb{R}^J \rightarrow \mathbb{R}^J$  is  $L$ -Lipschitz continuous, with respect to the two-norm on  $\mathbb{R}^J$ , if, for every  $x$  and  $y$  in  $\mathbb{R}^J$  we have

$$(1.1) \quad \|Tx - Ty\|_2 \leq L\|x - y\|_2.$$

**Definition 1.2.** An operator  $N$  on  $\mathbb{R}^J$  is called non-expansive (ne), with respect to the two-norm on  $\mathbb{R}^J$ , if, for every  $x$  and  $y$  in  $\mathbb{R}^J$ , we have

$$(1.2) \quad \|x - y\|_2 \geq \|Nx - Ny\|_2.$$

Clearly, if  $T$  is  $L$ -Lipschitz continuous, then the operator  $N = \frac{1}{L}T$  is non-expansive.

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**Definition 1.3.** An operator  $F$  on  $\mathbb{R}^J$  is called firmly non-expansive (fne) if, for every  $x$  and  $y$  in  $\mathbb{R}^J$ , we have

$$(1.3) \quad \langle Fx - Fy, x - y \rangle \geq \|Fx - Fy\|_2^2.$$

Using Cauchy's Inequality, we show easily that a firmly non-expansive operator on  $\mathbb{R}^J$  is non-expansive. The following lemma provides a useful characterization of fne operators.

**Lemma 1.4.** *An operator  $F : \mathbb{R}^J \rightarrow \mathbb{R}^J$  is fne if and only if  $F = \frac{1}{2}(I + N)$ , for some operator  $N$  that is ne with respect to the two-norm.*

*Proof.* Suppose that  $F = \frac{1}{2}(I + N)$ . We show that  $F$  is fne if and only if  $N$  is ne in the two-norm. First, we have

$$\langle Fx - Fy, x - y \rangle = \frac{1}{2}\|x - y\|_2^2 + \frac{1}{2}\langle Nx - Ny, x - y \rangle.$$

Also,

$$\left\| \frac{1}{2}(I + N)x - \frac{1}{2}(I + N)y \right\|_2^2 = \frac{1}{4}\|x - y\|_2^2 + \frac{1}{4}\|Nx - Ny\|_2^2 + \frac{1}{2}\langle Nx - Ny, x - y \rangle.$$

Therefore,

$$\langle Fx - Fy, x - y \rangle \geq \|Fx - Fy\|_2^2$$

if and only if

$$\|Nx - Ny\|_2^2 \leq \|x - y\|_2^2.$$

□

**Corollary 1.5.** *An operator  $F$  is fne if and only if  $I - F$  is fne.*

**Definition 1.6.** An operator  $A : \mathbb{R}^J \rightarrow \mathbb{R}^J$  is averaged (av) if  $A = (1 - \alpha)I + \alpha N$ , for some operator  $N$  that is non-expansive with respect to the two-norm, and some scalar  $\alpha$  in  $[0, 1)$ .

It is clear from the definitions and Lemma 1.4 that any fne operator is av. The product of finitely many av operators is again av [1]. According to the Krasnosel'skii-Mann Theorem [18, 16], if  $A$  is averaged and has fixed points, then the sequence  $\{A^n x^0\}$  converges to a fixed point of  $A$ , for every initial vector  $x^0$ . The following theorem is well known; see, for example, [15, 19].

**Theorem 1.7.** *Let  $f(x)$  be convex and differentiable and its derivative,  $\nabla f(x)$ , non-expansive in the two-norm. Then  $\nabla f(x)$  is firmly non-expansive.*

Suppose that  $g(x) : \mathbb{R}^J \rightarrow \mathbb{R}$  is convex and the function  $\nabla g(x)$  is  $L$ -Lipschitz continuous. Let  $f(x) = \frac{1}{L}g(x)$ , so that  $\nabla f(x)$  is a non-expansive operator. According to Theorem 1.7, the operator  $\nabla f(x) = \frac{1}{L}\nabla g(x)$  is firmly non-expansive. The proof of Theorem 1.7 is not trivial. In [15] Golshtein and Tretyakov prove the following theorem, from which Theorem 1.7 follows immediately.

**Theorem 1.8.** *Let  $g : \mathbb{R}^J \rightarrow \mathbb{R}$  be convex and differentiable. The following are equivalent:*

1)

$$(1.4) \quad \|\nabla g(x) - \nabla g(y)\|_2 \leq \|x - y\|_2;$$

2)

$$(1.5) \quad g(x) \geq g(y) + \langle \nabla g(y), x - y \rangle + \frac{1}{2} \|\nabla g(x) - \nabla g(y)\|_2^2;$$

and

3)

$$(1.6) \quad \langle \nabla g(x) - \nabla g(y), x - y \rangle \geq \|\nabla g(x) - \nabla g(y)\|_2^2.$$

The proof of Theorem 1.8 given in [15] is repeated in [8].

If  $f : \mathbb{R}^J \rightarrow \mathbb{R}$  is convex and differentiable, and  $\nabla f$  is  $L$ -Lipschitz continuous, then  $A = I - \gamma \nabla f$  is averaged, for any  $\gamma$  in the interval  $(0, 2/L)$ .

**1.2. Moreau's Proximity Operators.** Let  $f : \mathbb{R}^J \rightarrow \mathbb{R}$  be convex. For each  $z \in \mathbb{R}^J$  the function

$$(1.7) \quad m_f(z) := \min_x \left\{ f(x) + \frac{1}{2} \|x - z\|_2^2 \right\}$$

is minimized by a unique  $x$  [24]. The operator that associates with each  $z$  the minimizing  $x$  is Moreau's proximity operator, and we write  $x = \text{prox}_f(z)$ . The operator  $\text{prox}_f$  extends the notion of orthogonal projection onto a closed convex set [20, 21, 22]. We have  $x = \text{prox}_f(z)$  if and only if  $z - x \in \partial f(x)$ , where the set  $\partial f(x)$  is the sub-differential of  $f$  at  $x$ , given by

$$(1.8) \quad \partial f(x) := \{u \mid \langle u, y - x \rangle \leq f(y) - f(x), \text{ for all } y\}.$$

Proximity operators are also firmly non-expansive [11]; indeed, the proximity operator  $\text{prox}_f$  is the resolvent of the maximal monotone operator  $B(x) = \partial f(x)$  and all such resolvent operators are firmly non-expansive [4].

**1.3. The Forward-Backward Splitting Algorithm.** Our objective here is to provide an elementary proof of convergence for the forward-backward splitting (FBS) algorithm; a detailed discussion of this algorithm and its history is given by Combettes and Wajs in [11].

Let  $f : \mathbb{R}^J \rightarrow \mathbb{R}$  be convex, with  $f = f_1 + f_2$ , both convex,  $f_2$  differentiable, and  $\nabla f_2$   $L$ -Lipschitz continuous. The iterative step of the FBS algorithm is

$$(1.9) \quad x^k = \text{prox}_{\gamma f_1} \left( x^{k-1} - \gamma \nabla f_2(x^{k-1}) \right).$$

The FBS iteration has the form  $x^k = Ax^{k-1}$ . When we select  $\gamma$  in the interval  $(0, 2/L)$ , the operator  $A$  becomes averaged, since it is now the product of a *fne* operator and an *av* operator. Convergence of the sequence  $\{x^k\}$  to a fixed point of  $A$ , whenever  $A$  has fixed points, then follows from the KM Theorem.

As we shall show, convergence of the sequence  $\{x^k\}$  to a fixed point of  $A$  can be established without using the KM Theorem or the machinery of *fne* and *av* operators, if  $\gamma$  is chosen to lie within the interval  $(0, 1/L]$ .

## 2. SEQUENTIAL UNCONSTRAINED OPTIMIZATION

Sequential unconstrained optimization algorithms can be used to minimize a function  $f : \mathbb{R}^J \rightarrow (-\infty, \infty]$  over a (not necessarily proper) subset  $C$  of  $\mathbb{R}^J$  [13]. At the  $k$ th step of a sequential unconstrained minimization method we obtain  $x^k$  by minimizing the function

$$(2.1) \quad G_k(x) = f(x) + g_k(x),$$

where the auxiliary function  $g_k(x)$  is appropriately chosen. If  $C$  is a proper subset of  $\mathbb{R}^J$  we may force  $g_k(x) = +\infty$  for  $x$  not in  $C$ , as in the barrier-function methods; then each  $x^k$  will lie in  $C$ . The objective is then to select the  $g_k(x)$  so that the sequence  $\{x^k\}$  converges to a solution of the problem, or failing that, at least to have the sequence  $\{f(x^k)\}$  converging to the infimum of  $f(x)$  over  $x$  in  $C$ .

Our main focus in this paper is the use of sequential unconstrained optimization algorithms to obtain iterative methods in which each iterate can be obtained in closed form. Now the auxiliary functions  $g_k(x)$  are selected not to impose a constraint, but to facilitate computation.

## 3. SUMMA

In [7] we presented a particular class of sequential unconstrained minimization methods called SUMMA. As we showed in that paper, this class is broad enough to contain barrier-function methods, proximal minimization methods, and the simultaneous multiplicative algebraic reconstruction technique (SMART). By reformulating the problem, the penalty-function methods can also be shown to be members of the SUMMA class. Any alternating minimization (AM) problem with the five-point property [12] can be reformulated as a SUMMA problem; therefore the expectation maximization maximum likelihood (EMML) algorithm for Poisson data, which is such an AM algorithm, must also be a SUMMA algorithm.

For a method to be in the SUMMA class we require that each auxiliary function  $g_k(x)$  satisfy the inequality

$$(3.1) \quad 0 \leq g_k(x) \leq G_{k-1}(x) - G_{k-1}(x^{k-1}),$$

for all  $x$ . Note that it follows that  $g_k(x^{k-1}) = 0$ , for all  $k$ .

We assume that the inequality in (3.1) holds for each  $k$ . We also assume that  $\inf f(x) = b > -\infty$ . The next two results are taken from [7].

**Proposition 3.1.** *The sequence  $\{f(x^k)\}$  is non-increasing and the sequence  $\{g_k(x^k)\}$  converges to zero.*

*Proof.* We have

$$f(x^{k+1}) + g_{k+1}(x^{k+1}) = G_{k+1}(x^{k+1}) \leq G_{k+1}(x^k) = f(x^k).$$

Therefore,

$$f(x^k) - f(x^{k+1}) \geq g_{k+1}(x^{k+1}).$$

The sequence  $\{f(x^k)\}$  is decreasing to a finite limit, since it is bounded below by  $b$ , and, therefore, the sequence  $\{g_k(x^k)\}$  converges to zero.  $\square$

**Theorem 3.2.** *The sequence  $\{f(x^k)\}$  converges to  $b$ .*

*Proof.* Suppose that there is  $\delta > 0$  such that  $f(x^k) \geq b + 2\delta$ , for all  $k$ . Then there is  $z \in C$  such that  $f(x^k) \geq f(z) + \delta$ , for all  $k$ . From the inequality in (3.1) we have

$$(3.2) \quad g_k(z) - g_{k+1}(z) \geq f(x^k) + g_k(x^k) - f(z) \geq f(x^k) - f(z) \geq \delta,$$

for all  $k$ . But this cannot happen; the successive differences of a non-increasing sequence of non-negative terms must converge to zero.  $\square$

#### 4. CONVERGENCE OF THE FBS ALGORITHM

Let  $f : \mathbb{R}^J \rightarrow \mathbb{R}$  be convex, with  $f = f_1 + f_2$ , both convex,  $f_2$  differentiable, and  $\nabla f_2$   $L$ -Lipschitz continuous. Let  $\{x^k\}$  be defined by Equation (1.9) and let  $0 < \gamma \leq 1/L$ .

For each  $k = 1, 2, \dots$  let

$$(4.1) \quad G_k(x) = f(x) + \frac{1}{2\gamma} \|x - x^{k-1}\|_2^2 - D_{f_2}(x, x^{k-1}),$$

where

$$(4.2) \quad D_{f_2}(x, x^{k-1}) = f_2(x) - f_2(x^{k-1}) - \langle \nabla f_2(x^{k-1}), x - x^{k-1} \rangle.$$

Since  $f_2(x)$  is convex,  $D_{f_2}(x, y) \geq 0$  for all  $x$  and  $y$  and is the Bregman distance formed from the function  $f_2$  [3].

The auxiliary function

$$(4.3) \quad g_k(x) = \frac{1}{2\gamma} \|x - x^{k-1}\|_2^2 - D_{f_2}(x, x^{k-1})$$

can be rewritten as

$$(4.4) \quad g_k(x) = D_h(x, x^{k-1}),$$

where

$$(4.5) \quad h(x) = \frac{1}{2\gamma} \|x\|_2^2 - f_2(x).$$

Therefore,  $g_k(x) \geq 0$  whenever  $h(x)$  is a convex function.

We know that  $h(x)$  is convex if and only if

$$(4.6) \quad \langle \nabla h(x) - \nabla h(y), x - y \rangle \geq 0,$$

for all  $x$  and  $y$ . This is equivalent to

$$(4.7) \quad \frac{1}{\gamma} \|x - y\|_2^2 - \langle \nabla f_2(x) - \nabla f_2(y), x - y \rangle \geq 0.$$

Since  $\nabla f_2$  is  $L$ -Lipschitz, the inequality (4.7) holds for  $0 < \gamma \leq 1/L$ .

**Lemma 4.1.** *The  $x^k$  that minimizes  $G_k(x)$  over  $x$  is given by Equation (1.9).*

*Proof.* We know that  $x^k$  minimizes  $G_k(x)$  if and only if

$$0 \in \nabla f_2(x^k) + \frac{1}{\gamma}(x^k - x^{k-1}) - \nabla f_2(x^k) + \nabla f_2(x^{k-1}) + \partial f_1(x^k),$$

or, equivalently,

$$\left( x^{k-1} - \gamma \nabla f_2(x^{k-1}) \right) - x^k \in \partial(\gamma f_1)(x^k).$$

Consequently,

$$x^k = \text{prox}_{\gamma f_1}(x^{k-1} - \gamma \nabla f_2(x^{k-1})).$$

□

**Theorem 4.2.** *The sequence  $\{x^k\}$  converges to a minimizer of the function  $f(x)$ , whenever minimizers exist.*

*Proof.* A relatively simple calculation shows that

$$(4.8) \quad G_k(x) - G_k(x^k) = \frac{1}{2\gamma} \|x - x^k\|_2^2 + \left( f_1(x) - f_1(x^k) - \frac{1}{\gamma} \langle (x^{k-1} - \gamma \nabla f_2(x^{k-1})) - x^k, x - x^k \rangle \right).$$

Since

$$(x^{k-1} - \gamma \nabla f_2(x^{k-1})) - x^k \in \partial(\gamma f_1)(x^k),$$

it follows that

$$\left( f_1(x) - f_1(x^k) - \frac{1}{\gamma} \langle (x^{k-1} - \gamma \nabla f_2(x^{k-1})) - x^k, x - x^k \rangle \right) \geq 0.$$

Therefore,

$$(4.9) \quad G_k(x) - G_k(x^k) \geq \frac{1}{2\gamma} \|x - x^k\|_2^2 \geq g_{k+1}(x).$$

Therefore, the inequality in (3.1) holds and the iteration fits into the SUMMA class.

Now let  $\hat{x}$  minimize  $f(x)$  over all  $x$ . Then

$$\begin{aligned} G_k(\hat{x}) - G_k(x^k) &= f(\hat{x}) + g_k(\hat{x}) - f(x^k) - g_k(x^k) \\ &\leq f(\hat{x}) + G_{k-1}(\hat{x}) - G_{k-1}(x^{k-1}) - f(x^k) - g_k(x^k), \end{aligned}$$

so that

$$\left( G_{k-1}(\hat{x}) - G_{k-1}(x^{k-1}) \right) - \left( G_k(\hat{x}) - G_k(x^k) \right) \geq f(x^k) - f(\hat{x}) + g_k(x^k) \geq 0.$$

Therefore, the sequence  $\{G_k(\hat{x}) - G_k(x^k)\}$  is decreasing and the sequences  $\{g_k(x^k)\}$  and  $\{f(x^k) - f(\hat{x})\}$  converge to zero.

From

$$G_k(\hat{x}) - G_k(x^k) \geq \frac{1}{2\gamma} \|\hat{x} - x^k\|_2^2,$$

it follows that the sequence  $\{x^k\}$  is bounded and that a subsequence converges to some  $x^*$  with  $f(x^*) = f(\hat{x})$ .

Replacing the generic  $\hat{x}$  with  $x^*$ , we find that  $\{G_k(x^*) - G_k(x^k)\}$  is decreasing, and by Equation (4.8), a subsequence, and therefore, the entire sequence, converges to zero. From the inequality in (4.9), we conclude that the sequence  $\{\|x^* - x^k\|_2^2\}$  converges to zero, and so  $\{x^k\}$  converges to  $x^*$ . This completes the proof of the theorem. □

### 5. SOME EXAMPLES

We present some examples to illustrate the application of the convergence theorem.

**5.1. Projected Gradient Descent.** Let  $C$  be a non-empty, closed convex subset of  $\mathbb{R}^J$  and  $f_1(x) = \iota_C(x)$ , the function that is  $+\infty$  for  $x$  not in  $C$  and zero for  $x$  in  $C$ . Then  $\iota_C(x)$  is convex, but not differentiable. We have  $\text{prox}_{\gamma f_1} = P_C$ , the orthogonal projection onto  $C$ . The iteration in Equation (1.9) becomes

$$(5.1) \quad x^k = P_C\left(x^{k-1} - \gamma \nabla f_2(x^{k-1})\right).$$

The sequence  $\{x^k\}$  converges to a minimizer of  $f_2$  over  $x \in C$ , whenever such minimizers exist, for  $0 < \gamma \leq 1/L$ .

**5.1.1. The CQ Algorithm.** Let  $A$  be a real  $I$  by  $J$  matrix,  $C \subseteq \mathbb{R}^J$ , and  $Q \subseteq \mathbb{R}^I$ , both closed convex sets. The split feasibility problem (SFP) is to find  $x$  in  $C$  such that  $Ax$  is in  $Q$ . The function

$$(5.2) \quad f_2(x) = \frac{1}{2} \|P_Q Ax - Ax\|_2^2$$

is convex, differentiable and  $\nabla f_2$  is  $L$ -Lipschitz for  $L = \rho(A^T A)$ , the spectral radius of  $A^T A$ . The gradient of  $f_2$  is

$$(5.3) \quad \nabla f_2(x) = A^T (I - P_Q) Ax.$$

We want to minimize the function  $f_2(x)$  over  $x$  in  $C$ , or, equivalently, to minimize the function  $f(x) = \iota_C(x) + f_2(x)$ . The projected gradient descent algorithm has the iterative step

$$(5.4) \quad x^k = P_C\left(x^{k-1} - \gamma A^T (I - P_Q) Ax^{k-1}\right);$$

this iterative method was called the *CQ*-algorithm in [5, 6]. The sequence  $\{x^k\}$  converges to a solution whenever  $f_2$  has a minimum on the set  $C$ , for  $0 < \gamma \leq 1/L$ .

In [10, 9] the *CQ* algorithm was extended to a multiple-sets algorithm and applied to the design of protocols for intensity-modulated radiation therapy.

**5.1.2. The Projected Landweber Algorithm.** The problem is to minimize the function

$$f_2(x) = \frac{1}{2} \|Ax - b\|_2^2,$$

over  $x \in C$ . This is a special case of the SFP and we can use the *CQ*-algorithm, with  $Q = \{b\}$ . The resulting iteration is the projected Landweber algorithm [2]; when  $C = \mathbb{R}^J$  it becomes the Landweber algorithm [17].

## 6. MINIMIZING $f_2$ OVER A LINEAR MANIFOLD

Suppose that we want to minimize  $f_2$  over  $x$  in the linear manifold  $M = S + p$ , where  $S$  is a subspace of  $\mathbb{R}^J$  of dimension  $I < J$  and  $p$  is a fixed vector. Let  $A$  be an  $I$  by  $J$  matrix such that the  $I$  columns of  $A^T$  form a basis for  $S$ . For each  $z \in \mathbb{R}^I$  let

$$d(z) = f_2(A^T z + p),$$

so that  $d$  is convex, differentiable, and its gradient,

$$\nabla d(z) = A \nabla f_2(A^T z + p),$$

is  $K$ -Lipschitz continuous, for  $K = \rho(A^T A)L$ . The sequence  $\{z^k\}$  defined by

$$(6.1) \quad z^k = z^{k-1} - \gamma \nabla d(z^{k-1})$$

converges to a minimizer of  $d$  over all  $z$  in  $\mathbb{R}^I$ , whenever minimizers exist, for  $0 < \gamma \leq 1/K$ .

From Equation (6.1) we get

$$(6.2) \quad x^k = x^{k-1} - \gamma A^T A \nabla f_2(x^{k-1}),$$

with  $x^k = A^T z^k + p$ . The sequence  $\{x^k\}$  converges to a minimizer of  $f_2$  over all  $x$  in  $M$ .

Suppose now that we begin with an algorithm having the iterative step

$$(6.3) \quad x^k = x^{k-1} - \gamma A^T A \nabla f_2(x^{k-1}),$$

where  $A$  is any real  $I$  by  $J$  matrix having rank  $I$ . Let  $x^0$  be in the range of  $A^T$ , so that  $x^0 = A^T z^0$ , for some  $z^0 \in \mathbb{R}^I$ . Then each  $x^k = A^T z^k$  is again in the range of  $A^T$ , and we have

$$(6.4) \quad A^T z^k = A^T z^{k-1} - \gamma A^T A \nabla f_2(A^T z^{k-1}).$$

With  $d(z) = f_2(A^T z)$ , we can write Equation (6.4) as

$$(6.5) \quad A^T (z^k - (z^{k-1} - \gamma \nabla d(z^{k-1}))) = 0.$$

Since  $A$  has rank  $I$ ,  $A^T$  is one-to-one, so that

$$(6.6) \quad z^k - z^{k-1} - \gamma \nabla d(z^{k-1}) = 0.$$

The sequence  $\{z^k\}$  converges to a minimizer of  $d$ , over all  $z \in \mathbb{R}^I$ , whenever such minimizers exist, for  $0 < \gamma \leq 1/K$ . Therefore, the sequence  $\{x^k\}$  converges to a minimizer of  $f_2$  over all  $x$  in the range of  $A^T$ .

## 7. FEASIBLE-POINT ALGORITHMS

Suppose that we want to minimize a convex differentiable function  $f(x)$  over  $x$  such that  $Ax = b$ , where  $A$  is an  $I$  by  $J$  full-rank matrix, with  $I < J$ . If  $Ax^k = b$  for each of the vectors  $\{x^k\}$  generated by the iterative algorithm, we say that the algorithm is a feasible-point method.

**7.1. The Projected Gradient Algorithm.** Let  $C$  be the feasible set of all  $x$  in  $\mathbb{R}^J$  such that  $Ax = b$ . For every  $z$  in  $\mathbb{R}^J$ , we have

$$(7.1) \quad P_C z = P_{NS(A)} z + A^T (AA^T)^{-1} b,$$

where  $NS(A)$  is the null space of  $A$ . Using

$$(7.2) \quad P_{NS(A)} z = z - A^T (AA^T)^{-1} Az,$$

we have

$$(7.3) \quad P_C z = z + A^T (AA^T)^{-1} (b - Az).$$

Using Equation (1.9), we get the iteration step for the projected gradient algorithm:

$$(7.4) \quad x^k = x^{k-1} - \gamma P_{NS(A)} \nabla f(x^{k-1}),$$

which converges to a solution for  $0 < \gamma \leq 1/L$ , whenever solutions exist.



In the next subsection we present a somewhat simpler approach.

**7.2. The Reduced Gradient Algorithm.** Let  $x^0$  be a feasible point, that is,  $Ax^0 = b$ . Then  $x = x^0 + p$  is also feasible if  $p$  is in the null space of  $A$ , that is,  $Ap = 0$ . Let  $Z$  be a  $J$  by  $J - I$  matrix whose columns form a basis for the null space of  $A$ . We want  $p = Zv$  for some  $v$ . The best  $v$  will be the one for which the function

$$\phi(v) = f(x^0 + Zv)$$

is minimized. We can apply to the function  $\phi(v)$  the steepest descent method, or the Newton-Raphson method, or any other minimization technique.

The steepest descent method, applied to  $\phi(v)$ , is called the reduced steepest descent algorithm [23]. The gradient of  $\phi(v)$ , also called the reduced gradient, is

$$\nabla\phi(v) = Z^T\nabla f(x),$$

where  $x = x^0 + Zv$ ; the gradient operator  $\nabla\phi$  is then  $K$ -Lipschitz, for  $K = \rho(A^T A)L$ .

Let  $x^0$  be feasible. The iteration in Equation (1.9) now becomes

$$(7.5) \quad v^k = v^{k-1} - \gamma\nabla\phi(v^{k-1}),$$

so that the iteration for  $x^k = x^0 + Zv^k$  is

$$(7.6) \quad x^k = x^{k-1} - \gamma ZZ^T\nabla f(x^{k-1}).$$

The vectors  $x^k$  are feasible and the sequence  $\{x^k\}$  converges to a solution, whenever solutions exist, for any  $0 < \gamma < \frac{1}{K}$ .

**7.3. The Reduced Newton-Raphson Method.** The same idea can be applied to the Newton-Raphson method. The Newton-Raphson method, applied to  $\phi(v)$ , is called the reduced Newton-Raphson method [23]. The Hessian matrix of  $\phi(v)$ , also called the reduced Hessian matrix, is

$$\nabla^2\phi(v) = Z^T\nabla^2 f(c)Z,$$

so that the reduced Newton-Raphson iteration becomes

$$(7.7) \quad x^k = x^{k-1} - Z\left(Z^T\nabla^2 f(x^{k-1})Z\right)^{-1}Z^T\nabla f(x^{k-1}).$$

Let  $c^0$  be feasible. Then each  $x^k$  is feasible. The sequence  $\{x^k\}$  is not guaranteed to converge.

## 8. CONCLUSIONS

The forward-backward splitting algorithm can be formulated as a member of the SUMMA class of sequential unconstrained minimization algorithms. Convergence of the iterative sequence can then be established without relying on the machinery of firmly non-expansive and averaged operators. Examples are given to illustrate the usefulness of the forward-backward splitting algorithm.

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